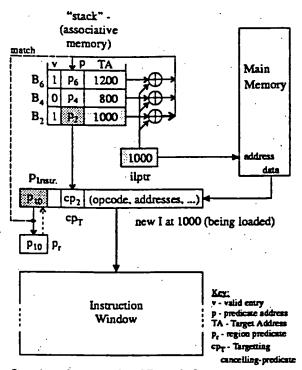


Disjoint

Nested

Overlapped

FIG. 1



Snapshot taken at t = 9+ of Example 5. - new I matches target address in stack

				predicate-ass (at load t		predicate-use (at code execution time)			
load <u>time</u>	address	code		Stack B v p TA	p _{in} =p _r cp _{in}	Pout	cp _{out} p _I	- condition for I execution	
1	100	\mathcal{I}_{1}	z = x op y	empty	. 1 0	p ₁ =1	•	1	
2 (<u> </u>] 200	B ₂	if (bc ₂) goto 400	B ₂ 1 P ₂ 400	1 0	$p_2 = \overline{bc_2}$	bc ₂	1	
		I ₃		B ₂ 1 P ₂ 400	p ₂ 0		÷	bc ₂	
3	400	I₄ ←		empxy	p ₂ cp ₂	bc ₂ +bc	·2 -	$\overline{bc}_2 + bc_2 = 1$	
5 🗓		I ₅		empty	. p ₄ 0	-	-	p ₄ =1	
6 🗐	600 #	B ₆	if (bc ₆) goto 800	B ₆ 1 P ₆ 800	p ₄ 0	bc ₆ ·p ₄	bc ₆ ·p ₄	1	
7	± 700	I,		B ₆ 1 P ₆ 800	p ₆ 0	•	-	bc ₆	
· 8].	800	I ₈ ←		empty	p ₆ cp ₆	bc ₆ +bc	c ₆ -	$\overline{bc}_6 + bc_6 = 1$	
9	900	. I ₉		етрху	p ₈ 0	-	•	p ₆ =1	
				Equations - fo	or "I": p _I =p _{out} =	p _{ia} +cp _{ia} ;	for "B": p _{ou}	= bc ·p _{in} , cp _{out} =bc·p _{in}	

丘6.3

				predicate (at lo	· · ·	predicate-use (at code execution time)			
load			•	stack			•		
<u>time</u>	address	code		<u>B</u> v p TA	p _{in} =p _r	<u>cp</u> in	Pout	cpout	p _I - condition for I execution
1	100	I ₁	z = x op y	empty	1	0	$p_1=1$	-	1
2	200	B ₂	if (bc ₂) goto 800	B ₂ 1 P ₂ 800	1	0	$p_2 = \overline{bc}_2$	bc ₂	1
3	300	I ₃		B ₂ 1 P ₂ 800	p ₂	.0	-	-	bc₂
4	400	B ₄ -	if (bc ₄) goto 600	B ₄ 1 P ₄ 600 B ₂ 1 P ₂ 800		0	bc4·b2	bc ₄ p ₂	1 .
5	500	I ₅		B ₄ 1 P ₄ 600 B ₂ 1 P ₂ 800	⊣ - ·	0	-	. -	$\overline{bc}_2 \cdot \overline{bc}_4$
6	600	I ₆ ←		B ₂ 1 P ₂ 800	p ₄	cp ₄	p ₄ +cp ₄	-	$\overline{bc_4} \cdot \overline{bc_2} + bc_4 \cdot \overline{bc_2} = \overline{bc_2}$
7	700	I ₇		B ₂ 1 P ₂ 800	P ₆	. 0	•	-	\overline{bc}_2
8	800	I ₈ ←		empty	P ₆	cp ₂	p ₆ +cp ₂	-	$\overline{bc}_2 + bc_2 = 1$
9	900	I ₉		етрху	P ₈	0	-	•	. 1
	900	19		етрку	P ₈	· · ·	<u> </u>		

 $Equations - \quad for \text{ ``I'': } p_{\overline{l}} = p_{out} = p_{in} + cp_{in}; \quad for \text{ ``B'': } p_{out} = \overline{bc} \cdot p_{in}, \quad cp_{out} = bc \cdot p_{in}$

FIG. 4

				predicate-as (at load				predicate-use (at code execution time)		
load				stack	•					
time	address	<u>code</u>		<u>B</u> v p TA	p _{in} =p _r	cp _{in}	<u> Pout</u>	cp _{out}	p ₁ - condition for I execution	
1	100	I ₁	z = x op y	етрху	1	0	$p_1=1$	- .'	1	
2	200	B ₂	if (bc ₂) goto 600	B ₂ 1 P ₂ 600	1	0	$p_2 = \overline{bc}_2$	bc ₂	1	
3	300	I ₃		B ₂ 1 P ₂ 600	P ₂	0	-	-	bc ₂	
4	400	B ₄	if (bc ₄) goto 800	B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 600	P ₂	0	bc ₄·p₂	bc ₄ ·p ₂	. 1	
5	500	I ₅		B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 600	P ₄	0	-	-	$\overline{bc_4} \cdot \overline{bc_2}$	
6	600	I ₆ ←		B ₄ 1 P ₄ 800 B ₂ 0 P ₂ 600	P4	cp ₂	p ₄ +cp ₂	-	$(\overline{bc}_4 \cdot \overline{bc}_2) + bc_2 = \overline{bc}_4 + bc_2$	
7	700	I,		B ₄ 1 P ₄ 800 B ₂ 0 P ₂ 600	P ₆	0	-	-	bc ₄+bc₂	
8	800	I ₈ ←		empty	P ₆	cp ₄	p ₆ +cp ₄	-	$\overline{bc}_4 + bc_2 + (bc_4 \cdot \overline{bc}_2) = 1$	
9	900	I ₉		empty	Pg	0	•	-	· 1	

Equations - for "T": $p_1 = p_{out} = p_{in} + cp_{in}$; for "B": $p_{out} = \overline{bc} \cdot p_{in}$, $cp_{out} = bc \cdot p_{in}$

FIG. 5

	· .				predicate-assignment (at load time)			٠.	predicate-use (at code execution time)		
	load	•			stack			•			
	<u>time</u>	address	<u>code</u>		B v p TA	$p_{in}=p_r$	<u>cp_{in}</u>	Pout	cp _{out}	p ₁ - condition for I execution	
	1	100	I ₁	z = x op y	empty	1	0	p ₁ =1	.	1 .	
•	2	200	B ₂ -	if (bc ₂) goto 1000	B ₂ 1 P ₂ 1000	1	0	p ₂ = bc ₂	bc ₂	1	
	3	300	I ₃		B ₂ 1 P ₂ 1000	P ₂	0		-	\overline{bc}_2	
	4 .	400	B ₄	if (bc ₄) goto 800	B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 1000	P ₂	0	bc ₄ ·p ₂	bc ₄ p ₂	1	
then the then the	5	500	I ₅		B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 1000	P ₄	0	•	-	bc₄·bc₂	
	6	600	В ₆	if (bc ₆) goto 1200	B ₆ 1 P ₆ 1200 B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 1000	P4	0	bc ₆ ⋅p ₄	bc ₆ p ₄	. 1	
	7	700	I ₇		B ₆ 1 P ₆ 1200 B ₄ 1 P ₄ 800 B ₂ 1 P ₂ 1000	P ₆	0	-	- ·	$\overline{bc}_6 \cdot \overline{bc}_4 \cdot \overline{bc}_2$	
	8	800	I _g ←		B ₆ 1 P ₆ 1200 B ₄ 0 P ₄ 800 B ₂ 1 P ₂ 1000	P ₆	ср ₄	p ₆ +cp ₄	-	$(\overline{bc}_{6} \overline{bc}_{4} \overline{bc}_{2}) + (bc_{4} \overline{bc}_{2})$ $= (\overline{bc}_{6} + bc_{4})\overline{bc}_{2}$	
	9	900	I ₉	·	B ₆ 1 P ₆ 1200 B ₄ 0 P ₄ 800 B ₂ 1 P ₂ 1000	p 8	0		-	(bc ₆ +bc ₄)bc ₂	
	- 10	1000	I ₁₀ ←		B ₆ 1 P ₆ 1200	P ₈	cp ₂	p ₈ +cp ₂	-	$((\overline{bc}_6 + bc_4)\overline{bc}_2) + bc_2$ $= \overline{bc}_6 + bc_4 + bc_2$	
	11	1100	I ₁₁	·	B ₆ 1 P ₆ 1200	p ₁₀	0	-	• -	$\overline{bc_6} + bc_4 + bc_2$	
	12	1200	I ₁₂ ←		empty	P 10	ср ₆	p ₁₀ +cp ₆	-	$ \overline{bc}_6 + bc_4 + bc_2 + (bc_6 \cdot \overline{bc}_4 \cdot \overline{bc}_2) $ =1	

empty

 $\text{Equations -} \quad \text{for "I": } p_{\overline{l}} = p_{out} = p_{in} + cp_{in}; \quad \text{ for "B": } p_{out} = \overline{bc} \cdot p_{in}, \quad cp_{out} = bc \cdot p_{in}$

FIG. 6

13

1300